

Linear equation systems: exact methods

Marcin Chrząszcz, Danny van Dyk
mchrzasz@cern.ch,
danny.van.dyk@gmail.com



**University of
Zurich** ^{UZH}

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Linear eq. system

⇒ This and the next lecture will focus on a well known problem. Solve the following equation system:

$$A \cdot x = b,$$

⇒ $A = a_{ij} \in \mathbb{R}^{n \times n}$ and $\det(A) \neq 0$

⇒ $b = b_i \in \mathbb{R}^n$.

⇒ The problem: Find the x vector.

Error digression

- ⇒ There is enormous amount of ways to solve the linear equation system.
- ⇒ The choice of one over the other of them should be gathered by the *condition* of the matrix A denoted at $cond(A)$. ⇒ If the $cond(A)$ is small we say that the problem is well conditioned, otherwise we say it's ill conditioned.
- ⇒ The *condition* relation is defined as:

$$cond(A) = \|A\| \cdot \|A^{-1}\|$$

- ⇒ Now there are many definitions of different norms... The most popular one (so-called "column norm"):

$$\|A\|_1 = \max_{1 \leq j \leq n} \sum_{i=1}^n |a_{i,j}|,$$

where n -is the dimension of A , i, j are columns and rows numbers.

More norms

⇒ A different norm is a spectral norm:

$$\|A\|_2 = \sqrt{\rho(A^T A)}$$

$$\rho(M) = \max\{|\lambda_i| : \det M - \lambda I = 0, i = 1, \dots, n\}$$

where $\rho(M)$ - spectral radius of M matrix, I unit matrix, λ_i eigenvalues of M .

⇒ Row norm:

$$\|A\|_\infty = \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{i,j}|,$$

Digression:

⇒ Calculation of the matrix norms are not a simple process at all. There are certain class of matrices that make the calculations easier.

⇒ The spectral norm can be also defined:

$$\text{cond}_2(A) = \frac{\max_{1 \leq i \leq n} |\lambda_i|}{\min_{1 \leq i \leq n} |\lambda_i|},$$

Example, ill-conditioned matrix

⇒ The text-book example of wrongly conditioned matrix is the Hilbert matrix:

$$h_{i,j} = \frac{1}{i+j-1}$$

⇒ Example:

$$h^{4 \times 4} = \begin{pmatrix} 1 & \frac{1}{2} & \frac{1}{3} & \frac{1}{4} \\ \frac{1}{2} & \frac{1}{3} & \frac{1}{4} & \frac{1}{5} \\ \frac{1}{3} & \frac{1}{4} & \frac{1}{5} & \frac{1}{6} \\ \frac{1}{4} & \frac{1}{5} & \frac{1}{6} & \frac{1}{7} \end{pmatrix}$$

⇒ The condition of this matrix:

$$\text{cond}(A) = \mathcal{O} \left(\frac{e^{3.52N}}{\sqrt{N}} \right)$$

⇒ For 8×8 matrix we get:

$$\text{cond}_1(A) = 3.387 \cdot 10^{10}, \quad \text{cond}_2(A) = 1.526 \cdot 10^{10}, \quad \text{cond}_\infty(A) = 3.387 \cdot 10^{10}$$

⇒ Clearly large numbers ;)

Exact methods: Cramer method

⇒ If $\det A \neq 0$ then the solutions are given by:

$$x_i = \frac{\det A_i}{\det A}$$

⇒ So calculate the solutions one needs to calculate $n + 1$ determinants. To calculate each determinate one needs $(n - 1)n!$ multiplications.

⇒ Putting it all together one needs $(n + 1)(n - 1)n! = n^{n+2}$

⇒ Brutal force but works ;)

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Exact methods: Gauss method

⇒ The idea besides the Gauss method is simple: transform the $Ax = b$ to get the equivalent matrix $A^{[n]}x = b^{[n]}$ where $A^{[n]}$ is triangular matrix:

$$A^{[n]} = \begin{pmatrix} a_{11}^{[n]} & a_{12}^{[n]} & \dots & a_{1n}^{[n]} \\ 0 & a_{22}^{[n]} & \dots & a_{2n}^{[n]} \\ \dots & & & \\ 0 & 0 & \dots & a_{nn}^{[n]} \end{pmatrix}$$

⇒ The algorithm: ⇒ To do so we calculate the: $d_{i,1}^{[1]} = \frac{a_{i1}^{[1]}}{a_{11}^{[1]}}$

⇒ The first row multiplied by the $d_{i,1}^{[1]}$ we subtract from the i^{th} row. ⇒ After this we get:

$$\begin{pmatrix} a_{11}^{[1]} & a_{12}^{[1]} & \dots & a_{1n}^{[1]} \\ 0 & a_{22}^{[1]} & \dots & a_{2n}^{[1]} \\ \dots & & & \\ 0 & a_{n2}^{[1]} & \dots & a_{nn}^{[1]} \end{pmatrix} \xrightarrow{x} = \begin{pmatrix} b_1^{[1]} \\ b_1^{[1]} \\ \dots \\ b_1^{[1]} \end{pmatrix}$$

Exact methods: Gauss method 2

⇒ Now one needs to repeat the above n times moving each time row down.

Cons:

- ⇒ The algorithm can be stopped if you divide by zero.
- ⇒ The method is very efficient to accumulate numerical errors.

Pros:

- ⇒ The number of needed floating point operations is less than Cramer.
- ⇒ Example for 15 equations: 1345 vs $5 \cdot 10^{12}$.

Backup