# Testing random number generators

Marcin Chrząszcz mchrzasz@cern.ch



University of Zurich<sup>uz</sup>

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# General methodology

### How to check if we have a good generator?

The generator is good if the number sequences that it produces have properties of truly random numbers. But how to check this!?!

 $\Rightarrow$  Traditional methodology:

Define some properties of random numbers from  $\mathcal{U}(0,1)$  and check if in the tests the properties are conserved.

- The problem with this approach is the fact that there are infinite number of test like this one would have to do :(
- In practice one can only only prove the generator is bad, but not that it's good.
- There is no way to guarantee that if the n tests are fulfilled the n + 1 will not fail!

⇒ The testing can be only in terms of so-called negative selection. ⇒ By each test our trust in the generator increases our trust in it, but it's not GM cars! There is no guarantee.

# General methodology, example

- $\Rightarrow$  Let's assume we have a generator that has  $\mathcal{U}(0,1)$ :
- We generate *n* numbers(*n* is fixed).
- From them we calculate a values of test function *T*.
- We calculate the F(T) where F is the CDF of the T statistics.
- Repeat the procedure N times:  $T_1, ..., T_N$  and  $F(T_1), ..., F(T_N)$ .

⇒ If the generator is good(hipothesis of  $\mathcal{U}(0,1)$  is true the  $F(T_1), ..., F(T_N)$  will have the distribution of  $\mathcal{U}(0,1)$ . One usually quotes the credibility level of a test! ⇒ There are number of test that the generator can be applied to. in the literature:

- "The Art of Computer Programming", Author Donald Knuth
- DIEHARD by G.Marsaglia, stat.fsu.edu/pub/diehard/

#### ⇒Modern approach:

Us the same formalism as is in studies the classical chaotic dynamical systems (same formalism is used in the modern generators).

 $\Rightarrow$  The RANLUX generator fulfils the chaotic test and all known classical tests (not surprisingly ;) )

# The $\chi^2$ texts with $\mathcal{U}(0,1)$

- $\Rightarrow$  The algorithm:
- Divide the [0,1) into k subdivisions:

$$0 = a_0 < a_1 < a_2 < a_3 < \dots < a_k = 1$$

- Let  $a_{n_i} = X_1, X_2, ... X_n$  be an series of elements in the interval  $[a_{i-1}, a_i)$  (with  $n_i$  elements). The  $p_i = P(a_{i-1} < X < a_i) = a_i a_{i-1}$ .
- A random variable:

$$\chi_k^2 = \sum_{i=1}^k \frac{(n_i - np_i)^2}{np_i}, \quad n = \sum_{i=1}^k n_i,$$

had a  $\chi^2$  distribution of k-1 degrees of freedom.

⇒ The above hypothesis verifies if the random numbers are indeed  $\mathcal{U}(0,1)$ . ⇒ The  $\chi^2$  distribution:  $X \in \mathbb{R}, X > 0, N \in \mathbb{N}$ :

$$\rho(X) = \frac{1}{2} \left(\frac{X}{2}\right)^{\frac{N}{2}-1} e^{\frac{X}{2}} \left[\Gamma\left(\frac{N}{2}\right)\right]^{-1} \quad E(X) = N, \quad V(X) = 2N$$

### The multi dimension test

 $\Rightarrow$  From the obtained numbers we construct an m dimension points:

$$(X_1, X_2, ..., X_m), (X_{m+1}, ..., X_{2m}), ..., (X_{(n-1)m+1}, X_{nm})$$

- In principle they should have a uniform distribution in an (0, 1)<sup>m</sup> hipercube.
- we divide each edge of the hipercube into k equal subdivisions: (j 1)/k, j/k), j = 1...k.
- Now: n<sub>i</sub> is the number of m dimensional points, which are in the i-th hipercube.
- The χ<sup>2</sup> test statistics:

$$\chi^2_{k^m-1} = \frac{k^m}{n} \sum_{i=1}^{k^m} n_i^2 - n, \quad n = \sum_{i=1}^{k^m} n_i$$

⇒ Now we construct other points:

$$(X_1, X_2, ..., X_m), (X_2, X_3, ..., X_{m+1}), \ldots$$

For N random numbers we have N − m + 1 such numbers.

We define the statistics:

$$\psi_0^2 = 0, \quad \psi_m^2 = \sum_{i=1}^{k^m} \frac{[n_i - (N - m + 1)/k^m]^2}{(N - m + 1)/k^m}, \quad m = 1, 2, ..$$

• For large N the  $(\psi_m^2-\psi_{m-1}^2)$  has a  $\chi^2$  distribution with  $k^m-k^{m-1}$  degrees of freedom.

<sup>D</sup>/15

# Overlaping-pairs-sparse-occupancy

 $\Rightarrow$  The OPSO (G.Marsaglia 1984)is an analysis of pairs obtained from random number generator.

 $X_1, X_2, ..., X_n$  - n random numbers obtained from generator. From each number we take b bits from which we construct a second series:  $I_1, I_2, ..., I_n$ , where  $I_j \in [0, 1, ..., 2^b - 1]$ .

 $\Rightarrow$  Next we create the pair series:

$$(I_1, I_2), (I_2, I_3), \dots (I_{n-1}, I_n)$$

 $\Rightarrow Y$  - number of pairs from :  $(i,j):i,j=0,1,...,2^b-1,$  which DIDN'T occur in the above series.

Bitstream	No. missing words	z-score	p-value
23 to 32	141989	0.2747	0.391764
22 to 31	142538	2.1678	0.015086
21 to 30	142084	0.6023	0.273484
20 to 29	142081	0.5920	0.276937

 $\Rightarrow$  This kind of test can be exteded to triple-pairs, and quadro-pairs.

⇒ See DIEHARD G.Marsgalia 1993 http://stat.fsu.edu/pub/diehard/

### Kolmogorov - Smirnov

 $\Rightarrow$  The K-S test is used to check if a Random variable has pdf of a distribution F. The test is based on the difference between the two distributions:

$$D_n = \sup_{-\infty < x < \infty} |F_n(x) - F(x)|, \quad F_n = \frac{1}{n} \sum_{j=1}^n \Theta(x - X_j).$$

 $\Rightarrow$  If the random generator is from the *F* distribution then the  $D_n \rightarrow 0$  with the probability 1.

 $\Rightarrow$  Large values of  $D_n$  exclude the generator.  $\Rightarrow$  The critical values of the test  $D_n(\alpha)$  can be find in the mathematical tables for every  $\alpha$ :

$$\mathcal{P}[D_n < D_n(\alpha)] = \alpha$$

⇒ They do not depend on the F function. ⇒ For the  $\mathcal{U}(0,1)$ :

$$F(x) = x, \ 0 < x < 1$$

# Kolmogorov - Smirnov in practice

#### Take note:

Empirical CDF of  $F_n$  is a step function and  $\sup_{-\infty < x < \infty} |F_n(x) - F(x)|$  is achieved only in one point!

 $\Rightarrow$  In practice one should sort the numbers:  $X_1, ..., X_n$  and calculate the following:

$$D_n^+ = \max_{1 \le i \le n} \left( \frac{i}{n} - F(X_{i:n}) \right), \quad D_n^- = \max_{1 \le i \le n} \left( F(X_{i:n}) - \frac{i-1}{n} \right)$$
$$D_n = \max\{D_n^+, D_n^-\}$$

where  $X_{i:n}$  is so-called position statistic:  $X_{1:n}, X_{2:n}, ..., X_{i:n}$ .  $\Rightarrow$  The statistic  $D_n$  asymptotically (in practice  $n \ge 80$ ) is approaching the  $\lambda$ -Kolomogorows cdf:

$$\lim_{n \to \infty} \mathcal{P}\{\sqrt{n}D_n \leqslant t\} = K(t) = \sum_{j=-\infty}^{\infty} (-1)^j e^{-2j^2t^2}, \ t > 0$$

for which the critical values  $\lambda_{\alpha}(\mathcal{P}\{\sqrt{n}D_n\}) > \lambda_{\alpha}$  can be found in the mathematical tables.

 $\Rightarrow$  Commonly the  $\lambda_{0.1} = 1.224$ ,  $\lambda_{0.05} = 1.358$ ,  $\lambda_{0.01} = 1.628$  are used.

### Statistic distributions test- sum test

 $\Rightarrow$  The *h* function has the form:

$$y = x_1 + x_2 + x_3 \dots x_m.$$

 $\Rightarrow$  the random variables form the new pdf:

$$g_m(y) = \begin{cases} \frac{1}{m-1} \left[ y^{m-1} - {m \choose 1} (y-1)^{m-1} + {m \choose 2} (y-2)^{m-1} - .. \right] & \text{for } 0 \leqslant y \leqslant m, \\ 0 & \text{else} \end{cases}$$

where you stop when y - m is negative.  $\Rightarrow$  For m = 2 we have the triangle pdf:

$$g_2(y) = \begin{cases} y, \text{ for } 0 \leq y \leq 1\\ 2-y, \text{ for } 0 \leq y \leq 1 \end{cases}$$

 $\Rightarrow$  For m = 3 we have the triangle pdf:

$$g_{3}(y) = \begin{cases} \frac{1}{2}y^{2}, \text{ for } 0 \leq y \leq 1\\ \frac{1}{2} \left[y^{2} - 3(y-1)^{2}\right], \text{ for } 1 \leq y \leq 2\\ \frac{1}{2} \left[y^{2} - 3(y-1)^{2} 3(y-2)^{2}\right], \text{ for } 2 \leq y \leq 3 \end{cases}$$

 $\Rightarrow$  For large *m* the  $g_m$  approaches the normal distribution.

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# Statistic distributions test- $d^2$

 $\Rightarrow$  for m = 4 we define the h:

$$y = (x_1 - x_3)^2 + (x_2 - X_4)^2$$

aka the square distance between  $(x_1, x_2)$  and  $(x_3, x_4)$ .  $\Rightarrow$  If the  $X_1, X_2, X_3, X_4$  are from  $\mathcal{U}(0, 1)$  then:

$$d^{2} = (X_{1} - X_{3})^{2} + (X_{2} - X_{4})^{2}$$

had a pdf given by the following formula:

$$\mathcal{P}(d^2 - y) = \begin{cases} \pi y - \frac{8}{3}y^{\frac{3}{2}} + \frac{1}{2}y^2 & \text{for } 0 \leq y \leq 1\\ -\frac{1}{2}y^2 - 4 \operatorname{arcsec}(y^{\frac{1}{2}}) & \text{for } 1 \leq y \leq 2 \end{cases}$$

 $\Rightarrow$  Test is to check if the generated numbers have the aforementioned distribution.

### Statistic distributions test- pair distance

 $\Rightarrow$  Generate n points from  $(0,1)^m.$  We take  $\binom{n}{2}$  pairs of points and we calculate the distance between them.

 $\Rightarrow$  If D is the smallest distance between the pairs  $\mapsto$  for the  $\mathcal{U}(0,1)^m$  the  $T = n^2 D^m/2$  has the exponential distribution with the mean  $1/V_m$ , where  $V_m$  is the hiper volume of the unite ball.

 $\Rightarrow$  In Patrice:

- We generate Nn points in the hipercube  $(0,1)^m$ , getting N points in the T statistics.
- We compare the empirical distribution *T* with the exponential distribution.
- WARNING: the N, n, m need to be choose smartly for the test to make sense.
- $\Rightarrow$  Linear generators usually fail this test!

### Statistic distributions test- series test

 $\Rightarrow$  Lets assume our numbers are generated with a CDF *F*. The values of *F* we divide in two separated sub-samples: *A* and *B*.

 $\Rightarrow$  Furthermore we define the new variables Y such as:

$$Y = \begin{cases} = aX \in A \\ = bX \in B \end{cases}$$

 $\Rightarrow$  The random number sequence we transform the  $X_1, X_2, X_3, ..., X_n$  into  $Y_1, Y_2, Y_3, ..., Y_n$ .

 $\Rightarrow$  Next we make series: For example the a, a, b, a, a, b, b, b, a will be grouped into aa, b, aa, bbb, a.

- $\Rightarrow$  Let  $n_a$  be number of a symbols in  $Y_1, Y_2, Y_3, ..., Y_n$ .  $n_b = N n_a$ .
- $\Rightarrow$  Distribution of number of series (*R*) is given by the equation:

$$\mathcal{P}(R=r,n_a,n_b) = \begin{cases} 2\binom{n_a-1}{k-1}\binom{n_b-1}{k-1} / \binom{N}{n_a} & \text{if } r=2k\\ [\binom{n_a-1}{k}\binom{n_b-1}{k-1} + \binom{n_a-1}{k-1}\binom{n_b-1}{k}] / \binom{N}{n_a} & \text{if } r=2k+1 \end{cases}$$

### Statistic distributions test- poker test

 $\Rightarrow$  The values of X random variable we divide into k identical sub samples:

 $0 < a_1 < \ldots < a_k = 1$ 

 $\Rightarrow$  For  $X_1, X_2, ..., X_n$  from  $\mathcal{U}(0, 1)$ :

$$\mathcal{P}(a_{i-1} < X_j < a_i) = \frac{1}{k}.$$

 $\Rightarrow$  We create the new variables  $Y_1$  accordingly:

$$Y_j = i$$
 if  $X_j \in (a_{i-1}, a_i), \ i = 0, 1, \dots k - 1$ 

 $\Rightarrow$  Now we create "the fives":

 $(Y_1, Y_2, Y_3, Y_4, Y_5), (Y_6, \dots$ 

 $\Rightarrow$  There are couple of types of fives:

aabcd pair

aaabc three

aaaab four

aaaaa five

# Statistic distributions test- poker test

 $\Rightarrow$  If the variables are independent then we can calculate the probability:

$$\begin{aligned} \mathcal{P}\{(abcde)\} &= \frac{(k-1)(k-2)(k-3)(k-4)}{k^4}, \quad k \ge 5, \\ \mathcal{P}\{(aabcd)\} &= \frac{10(k-1)(k-2)(k-3)}{k^4}, \quad k \ge 4, \\ \mathcal{P}\{(aabbc)\} &= \frac{15(k-1)(k-2)}{k^4}, \quad k \ge 3, \\ \mathcal{P}\{(aaabc)\} &= \frac{10(k-1)(k-2)}{k^4}, \quad k \ge 3, \\ \mathcal{P}\{(aaabb)\} &= \frac{10(k-1)}{k^4}, \quad k \ge 3, \\ \mathcal{P}\{(aaaab)\} &= \frac{5(k-1)}{k^4}, \quad k \ge 2, \\ \mathcal{P}\{(aaaaa)\} &= \frac{1}{k^4}, \quad k \ge 1, \end{aligned}$$

 $\Rightarrow$  In practice people choose: k = 2, 8, 10

 $\Rightarrow$  The agreemnt of the distribution of different types of fives is check using the  $\chi^2$  test.

Marcin Chrząszcz (Universität Zürich)

### Conclusions

 $\Rightarrow$  There are infinite number of tests one can invent for the testing of the generators.

 $\Rightarrow$  All of the tests are in the same taste: invent a problem where you know the analytic solution, solve the problem and compare the results.

 $\Rightarrow$  Homework: Use one of the previously implemented random number generator and :

- E5.1 Test them with chi-square test k=10.
- E5.2 Kolomorov-smirnon.
- E5.3 Multidimensional test.

# Backup